

Homework 6

Hint

Here is an example problem illustrating how to complete Homework 6.

Hint: Let V be the linear code over $GF(2)$ given by the following generator matrix

$$G = \begin{pmatrix} 1 & 0 & 0 & 0 & 1 & 1 & 0 \\ 0 & 1 & 0 & 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 0 & 1 & 1 & 1 \\ 0 & 0 & 0 & 1 & 1 & 0 & 1 \end{pmatrix}$$

- a) What is the length n of V ? What the dimension k of V ? Please explain your answer.

Since the generator matrix has 7 columns, $n=7$. Since the generator matrix is already in echelon canonical form, the dimension of V is equal to the number of non-zero rows. Hence, $k=4$.

- b) Find a parity check matrix H for V .

Since the generator matrix G is already in the form $G = (I | P)$, we know that the generator matrix H is $H = (-P^T | I)$. Hence,

$$H = \begin{pmatrix} 1 & 0 & 1 & 1 & 1 & 0 & 0 \\ 1 & 1 & 1 & 0 & 0 & 1 & 0 \\ 0 & 1 & 1 & 1 & 0 & 0 & 1 \end{pmatrix}$$

- c) Use the parity check matrix H to list all the elements of the dual code V^\perp .

Since $\text{Dim}(V^\perp) = n - k = 7 - 4 = 3$, it follows that the space of infowords for V^\perp is the space of all binary 3-tuples. Thus the matrix

$$W = \begin{pmatrix} 0 & 0 & 0 \\ 0 & 0 & 1 \\ 0 & 1 & 0 \\ 0 & 1 & 1 \\ 1 & 0 & 0 \\ 1 & 0 & 1 \\ 1 & 1 & 0 \\ 1 & 1 & 1 \end{pmatrix}$$

is a matrix whose rows are a complete list of the infowords of the code V^\perp . Since H is the generator matrix of V^\perp , it follows that rows of the matrix

$$W \cdot H = \begin{pmatrix} 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 1 & 1 & 0 & 0 & 1 \\ 1 & 1 & 1 & 0 & 0 & 1 & 0 \\ 1 & 0 & 0 & 1 & 0 & 1 & 1 \\ 1 & 0 & 1 & 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 & 1 & 0 & 1 \\ 0 & 1 & 0 & 1 & 1 & 1 & 0 \\ 0 & 0 & 1 & 0 & 1 & 1 & 1 \end{pmatrix}$$

are a complete list of the codewords of the code V^\perp .

Suggestion: Use Maple to compute this matrix. See:

- <http://www.cs.umbc.edu/~lomonaco/f06/653/hwk653/hwk06-hint-mapleworksheet.pdf>
- <http://www.cs.umbc.edu/~lomonaco/f06/653/hwk653/hwk06-hint.mw>)

d) Use the results of **c)** to create a weight distribution table, i.e., a table with two columns labeled respectively j and A_j^\perp .

j	A_j^\perp
0	1
1	0
2	0
3	0
4	7
5	0
6	0
7	0

e) Use the table created in **d)** to find the weight enumerator polynomial $A^\perp(x)$ of the dual code V^\perp .

$$A^\perp(x) = 1 + 7x^4$$

f) Use the MacWilliams identity to compute the weight enumerator polynomial $A(x)$ of the original linear code V .

MacWilliams identity for the weight enumerator $A(x)$ of V in terms of the weight enumerator $A^\perp(x)$ of V^\perp is

$$2^{n-k} A(x) = (1+x)^n A^\perp\left(\frac{1-x}{1+x}\right)$$

Thus,

$$\begin{aligned} A(x) &= 2^{-3} (1+x)^7 A^\perp\left(\frac{1-x}{1+x}\right) = \frac{1}{8} (1+x)^7 \left(1 + 7\left(\frac{1-x}{1+x}\right)^4\right) \\ &= \frac{1}{8} (1+x)^7 + \frac{7}{8} (1+x)^3 (1-x)^4 = 1 + 7x^3 + 7x^4 + x^7 \end{aligned}$$

Suggestion: Use Maple to compute polynomial. See:

- <http://www.cs.umbc.edu/~lomonaco/f06/653/hwk653/hwk06-hint-mapleworksheet.pdf>
- <http://www.cs.umbc.edu/~lomonaco/f06/653/hwk653/hwk06-hint.mw>)

g) Use the weight enumerator polynomial $A(x)$ to find the minimum distance of V .

Since the weight enumerator $A(x)$ of V is $A(x) = 1 + 7x^3 + 7x^4 + x^7$, we know that the code V consists of **16** vectors, **1** of weight **0**, **7** of weight **3**, **7** of weight **4**, and **1** of weight **7**. Thus, the minimum non-zero weight d of V is $d = 3$.