





CMSC 461, Database Management Systems Spring 2018

Lecture 11 - Chapter 8 Relational Database Design Part 1

These slides are based on "Database System Concepts" 6th edition book and are a modified version of the slides which accompany the book (http://codex.cs.yale.edu/avi/db-book/db6/slide-dir/index.html), in addition to the 2009/2012 CMSC 461 slides by Dr. Kalpakis

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https://www.csee.umbc.edu/~jsleem1/courses/461/spr18

Logistics

- HW3 released today due 3/12/2018
- Phase 2 due Wednesday 3/7/2018
- Midterm 3/14/2018

Lecture Outline

- Relational Database Design
- First Normal Form (1NF)
- Functional Dependencies
- Boyce-Codd (BCNF)
- Third Normal Form (intro)

Lecture Outline

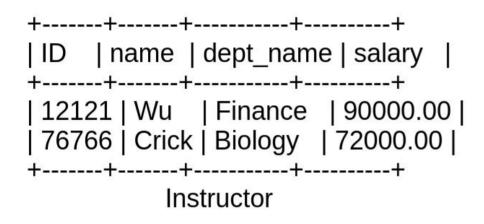
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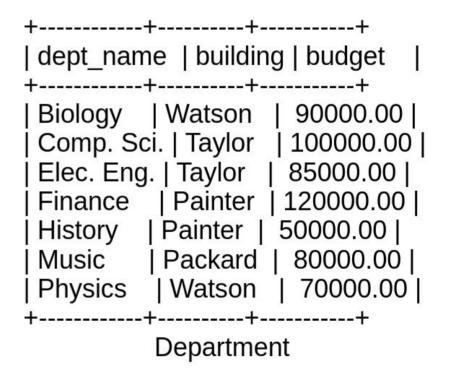
Relational Database Design

- We want to move from the E-R diagram to a set of relations
 - Eliminate redundancy
 - Ensure design is complete
 - Ensure information is easily retrievable
- . Going to learn about
 - normal form
 - functional dependencies

Design Alternatives: Combining Schemas

• Recall the instructor and department relations





Design Alternatives: Combining Schemas

- Suppose we combine instructor and department into inst_dept
 - No connection to relationship set inst_dept
- Why is this a bad idea?

ID	name	salary	dept_name	building	budget
22222	Einstein	95000	Physics	Watson	70000
12121	Wu	90000	Finance	Painter	120000
32343	El Said	60000	History	Painter	50000
45565	Katz	75000	Comp. Sci.	Taylor	100000
98345	Kim	80000	Elec. Eng.	Taylor	85000
76766	Crick	72000	Biology	Watson	90000
10101	Srinivasan	65000	Comp. Sci.	Taylor	100000
58583	Califieri	62000	History	Painter	50000
83821	Brandt	92000	Comp. Sci.	Taylor	100000
15151	Mozart	40000	Music	Packard	80000
33456	Gold	87000	Physics	Watson	70000
76543	Singh	80000	Finance	Painter	120000

Design Alternatives: Combining Schemas

Repetition of information

_					
ID	name	salary	dept_name	building	budget
22222	Einstein	95000	Physics	Watson	70000
12121	Wu	90000	Finance	Painter	120000
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Combining Schemas without Repetition?

• Consider combining relations:

sec_class(sec_id, building, room_number) and
section(course_id, sec_id, semester, year)

1 Painter 514	BIO-101	1	Summer	2009

• into one relation

section(course_id, sec_id, semester, year, building,
room_number)

course_id	sec_id	semester	year	building	room_number
BIO-101	1	Summer	2009	Painter	514

Given inst_dept was the result of our design.....

How would we know to split up (*decompose*) it into *instructor* and *department*?

• What can we say about this data?

ID	name	salary	dept_name	building	budget
22222	Einstein	95000	Physics	Watson	70000
12121	Wu	90000	Finance	Painter	120000
32343	El Said	60000	History	Painter	50000
45565	Katz	75000	Comp. Sci.	Taylor	100000
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• Write a rule:

"if there were a schema (*dept_name, building, budget*), then *dept_name* would be able to serve as the primary key"

- Denote as a *functional dependency*: *dept_name → building, budget*
- In *inst_dept*, because *dept_name* is not a primary key, the building and budget of a department may have to be repeated
 - This indicates the need to decompose *inst_dept*

Not all decompositions are good.

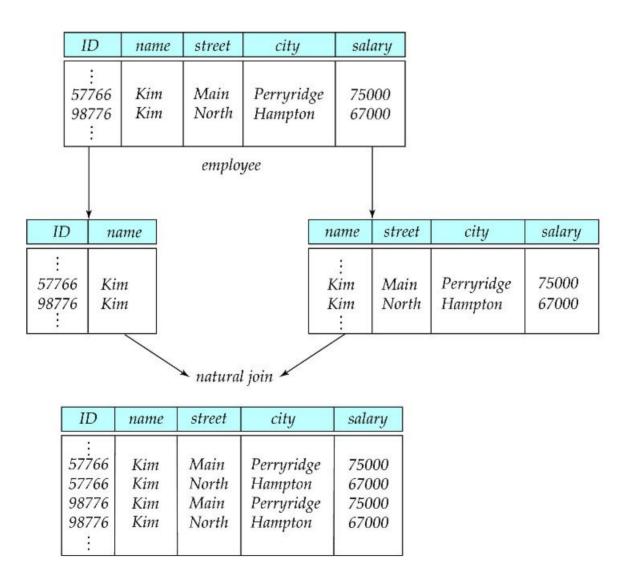
Suppose we decompose:

employee(ID, name, street, city, salary)

into

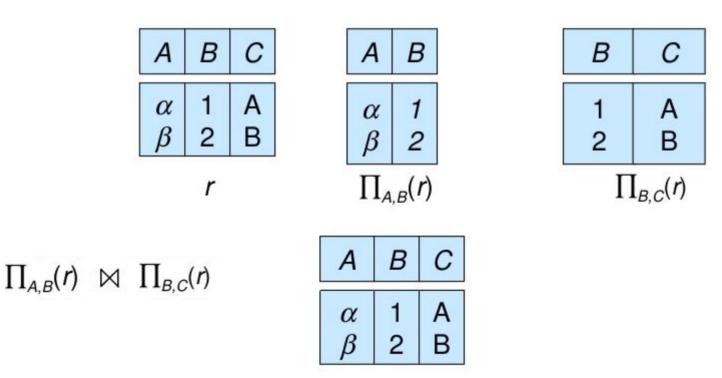
employee1 (ID, name) employee2 (name, street, city, salary)

A Lossy Decomposition



Lossless Join Decomposition

Decomposition of R = (A, B, C) $R_1 = (A, B)$ $R_2 = (B, C)$



Based on and image from "Database System Concepts" book and slides, 6th edition

Another Example

- Consider the relation schema:
 - Lending-schema = (branchName, branchCity, assets, customerName, loanNumber, amount)

branch-name	branch-city	assets	customer- name	loan- number	amount
Downtown	Brooklyn	9000000	Jones	L-17	1000
Redwood	Palo Alto	2100000	Smith	L-23	2000
Perryridge	Horseneck	1700000	Hayes	L-15	1500
Downtown	Brooklyn	9000000	Jackson	L-14	1500

Another Example

- Redundancy:
 - Data for branchName, branchCity, assets are repeated for each loan that a branch makes
 - Wastes space
 - Complicates updating, introducing possibility of inconsistency of assets value
- Null values
 - Cannot store information about a branch if no loans exist
 - Can use null values, but they are difficult to handle.

Decomposition

Decompose the relation schema Lending-schema into:

BranchSchema = (branchName, branchCity,assets)

LoanSchema = (customerName, loanNumber,branchName, amount)

All attributes of an original schema (R) must appear in the decomposition (R_1, R_2) :

$R = R_1 \cup R_2$

Lossless-join decomposition. For all possible relations *r* on schema *R*

$$r = \prod_{R_1} (r) \bowtie \prod_{R_2} (r)$$

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- Third Normal Form

- In the relational model attributes have no substructure
 - composite attributes each component becomes its own attribute
 - multivalued attributes one tuple for each item
- Domain is *atomic* if its elements are considered to be indivisible units
 - Examples of non-atomic domains:
 - Set of names, composite attributes
 - Identification numbers like CS001 where department is combined with employee number
 - Fine line: 'CS101' might be a course identifier and could be interpreted as atomic

- Non-atomic values complicate storage and encourage redundant (repeated) storage of data
 - Example: Set of accounts stored with each customer, and set of owners stored with each account

• A relational schema R is in *first normal form* if the domains of all attributes of R are atomic

- Atomicity is actually a property of how the elements of the domain are used
 - Example: Strings would normally be considered indivisible
 - Suppose that students are given roll numbers which are strings of the form CS0012 or EE1127
 - If the first two characters are extracted to find the department, the domain of roll numbers is not atomic.
 - Doing so is a bad idea: leads to encoding of information in application program rather than in the database

Is it atomic?

- Address
- Course ID (CS-101)
- Student Name
- SSN
- ISBN Number

Convert it to 1NF

ID	Color
1	Blue,Red
2	Yellow,Brown
3	Orange,Green

ID	Name	Address	Order
1	John Doe	13101 Brandley Lane	Cell Phone
2	Jackie Click	4531 Tinker Road	Charger
3	Brad Dunkin	8593 Gerwin Avenue	Cell Phone Case

Pitfalls in Relational Database Design

- Relational database design requires that we find a "good" collection of relation schemas
- A bad design may lead to
 - Repetition of Information
 - Inability to represent certain information
- Design Goals:
 - Avoid redundant data
 - Ensure that relationships among attributes are represented
 - Facilitate the checking of updates for violation of database integrity constraints

Goal: Devise a Theory for the following:

- Decide whether a particular relation R is in "good" form
- In the case that a relation R is not in "good" form, decompose it into a set of relations {R₁, R₂, ..., R_n } such that
 - each relation is in good form
 - the decomposition is a lossless-join decomposition
- This theory is based on:
 - functional dependencies
 - multivalued dependencies

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- Constraints on the set of legal relations
- Require that the value for a certain set of attributes determines uniquely the value for another set of attributes
- A functional dependency is a generalization of the notion of a key

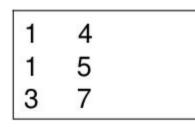
Remember

- *r* is a relation
- r(R) is the schema for the relation r
- *R* denotes the set of attributes
- K represents the set of attributes that is the superkey
- A superkey set of attributes that uniquely identify a tuple

- Let R be a relation schema $\alpha \subseteq R$ and $\beta \subseteq R$
- The functional dependency

 $\alpha \rightarrow \beta$

- holds on R if and only if for any legal relations r(R), whenever any two tuples t_1 and t_2 of r agree on the attributes α , they also agree on the attributes β . That is $t_1[\alpha] = t_2[\alpha] \implies t_1[\beta] = t_2[\beta]$
- Example: Consider r(A,B) with the following instance of r



 On this instance, A→B does NOT hold, but B→A does
 hold on and image from "Database System Concepts" book and slides, 6^{th edition}

 Functional dependencies allow us to express constraints that cannot be expressed using superkeys. Consider the schema:

inst_dept (ID, name, salary, dept_name, building, budget)

We expect these functional dependencies to hold: $dept_name \rightarrow building$ but would not expect the following to hold: $dept_name \rightarrow salary$

- We use functional dependencies to:
 - test relations to see if they are legal under a given set of functional dependencies
 - If a relation r is legal under a set F of functional dependencies, we say that r satisfies F
 - specify constraints on the set of legal relations
- We say that F holds on R if all legal relations on R satisfy the set of functional dependencies F
- Note: A specific instance of a relation schema may satisfy a functional dependency even if the functional dependency does not hold on all legal instances
 - For example, a specific instance of *instructor* may sometimes satisfy

- A functional dependency is trivial if it is satisfied by all instances of a relation
 - Example:
 - ID, name \rightarrow ID
 - $name \rightarrow name$
 - In general, $\alpha \rightarrow \beta$ is trivial if $\beta \subseteq \alpha$

Functional Dependencies Examples

Assume schema:

student(student_id, first_name, last_name, major, SSN)

Which are true in regards to functional dependencies:

 $student_id \rightarrow last_name$ $last_name \rightarrow student_id$ $student_id \rightarrow last_name, major, SSN, student_id$ $SSN \rightarrow student_id, last_name, major, SSN$ $first_name \rightarrow last_name$ $last_name \rightarrow last_name$

Functional Dependencies Examples

Assume schema:

student(student_id, first_name, last_name, major, SSN)

Which are true in regards to functional dependencies:

Closure of a Set of Functional Dependencies

- Given a set F of functional dependencies, there are certain other functional dependencies that are logically implied by F
 - For example:
 Given a schema r(A,B,C)
 If A → B and B → C
 then we can infer that A → C
- The set of all functional dependencies logically implied by F is the closure of F
- We denote the closure of *F* by *F*⁺ *F*⁺ is a superset of *F*

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Boyce-Codd Normal Form

A relation schema R is in BCNF with respect to a set F of functional dependencies if for all functional dependencies in F^+ of the form

 $\alpha \rightarrow \beta$

where $\alpha \subseteq R$ and $\beta \subseteq R$, at least one of the following holds:

- $\alpha \rightarrow \beta$ is trivial (i.e., $\beta \subseteq \alpha$)
- α is a superkey for *R*

Example schema *not* in BCNF:

instr_dept (ID, name, salary, dept_name, building, budget)

because *dept_name→ building*, *budget* holds on *instr_dept*, but *dept_name* is not a superkey

Boyce-Codd Normal Form

Are these schemas in BCNF:

instructor (<u>ID</u>, name, dept_name, salary) ID \rightarrow name,dept_name,salary

department(dept_name,building,budget) dept_name→ building, budget

Why or why not?

Based on and image from "Database System Concepts" book and slides, 6th edition

Boyce-Codd Normal Form

Are these schemas in BCNF:

instructor (<u>ID,</u>name, dept_name, salary) ID→ name,dept_name,salary

department(dept_name,building,budget) dept_name→ building, budget YES – ID is superkey

YES – dept_name is superkey

Decomposing a Schema into BCNF

• Suppose we have a schema *R* and a non-trivial dependency $\alpha \rightarrow \beta$ causes a violation of BCNF.

We decompose *R* into:

(α Uβ)

 $(R - (\beta - \alpha))$

• In our example:

instr_dept (<u>ID,</u> name, salary<u>, dept_name,</u> building, budget)

 $\alpha = dept_name$

 β = building, budget

and *inst_dept* is replaced by

 $(\alpha \cup \beta) = (dept_name, building, budget)$ $(R - (\beta - \alpha)) = (ID, name, salary, dept_name)$

Schema:

Student(ID,Name,AdvisorID,AdvisorName)

What are the functional dependencies?

Schema:

Student(ID,Name,AdvisorID,AdvisorName)

What are the functional dependencies?

ID -> Name AdvisorID -> AdvisorName

What uniquely identifies the tuples?

Schema:

Student(ID,Name,AdvisorID,AdvisorName)

What are the functional dependencies?

ID -> Name AdvisorID -> AdvisorName

What uniquely identifies the tuples? (ID,AdvisorID)

Is there a BCNF violation?

Schema:

Student(ID,Name,AdvisorID,AdvisorName)

What are the functional dependencies?

ID -> Name AdvisorID -> AdvisorName

What is the primary key? (ID,AdvisorID)

Is there a BCNF violation? YES!

Schema: Student(ID,Name,AdvisorID,AdvisorName)

What are the functional dependencies? ID -> Name AdvisorID -> AdvisorName

What is the primary key? (ID,AdvisorID)

Is there a BCNF violation? YES!

Use ID-> Name to decompose R- (Name-ID)= (ID,AdvisorID,AdvisorName) and ID union Name = (ID,Name)

BCNF and Dependency Preservation

- Constraints, including functional dependencies, are costly to check in practice unless they pertain to only one relation
- If it is sufficient to test only those dependencies on each individual relation of a decomposition in order to ensure that all functional dependencies hold, then that decomposition is dependency preserving
- Because it is not always possible to achieve both BCNF and dependency preservation, we consider a weaker normal form, known as third normal form

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Third Normal Form

 A relation schema R is in third normal form (3NF) if for all:

lpha
ightarrow eta in $F^{_+}$

at least one of the following holds:

- $\alpha \rightarrow \beta$ is trivial
- α is a superkey for *R*
- Each attribute A in $\beta \alpha$ is contained in a candidate key for R.

(NOTE: each attribute may be in a different candidate key)

- If a relation is in BCNF it is in 3NF (since in BCNF one of the first two conditions above must hold).
- Third condition is a minimal relaxation of BCNF to ensure dependency preservation (will see why later).