

1. Definitions of $\Theta(g(n))$, $O(g(n))$ and $\Omega(g(n))$.

$\Theta(g(n)) = \{ f(n) : \text{there exist positive constants } c_1, c_2, \text{ and } n_0 \text{ such that}$
 $0 \leq c_1 g(n) \leq f(n) \leq c_2 g(n) \text{ for all } n \geq n_0 \}.$

$O(g(n)) = \{ f(n) : \text{there exist positive constants } c \text{ and } n_0 \text{ such that}$
 $0 \leq f(n) \leq c g(n) \text{ for all } n \geq n_0 \}.$

$o(g(n)) = \{ f(n) : \text{for any positive constant } c, \text{ there exists a positive constant } n_0 \text{ such that}$
 $0 \leq f(n) < c g(n) \text{ for all } n \geq n_0 \}.$

$\Omega(g(n)) = \{ f(n) : \text{there exist positive constants } c \text{ and } n_0 \text{ such that}$
 $0 \leq c g(n) \leq f(n) \text{ for all } n \geq n_0 \}.$

$\omega(g(n)) = \{ f(n) : \text{for any positive constant } c, \text{ there exists a positive constant } n_0 \text{ such that}$
 $0 \leq c g(n) < f(n) \text{ for all } n \geq n_0 \}.$

2. Master Theorem.

Let $a \geq 1$ and $b > 1$ be constants, let $f(n)$ be a function, and let $T(n)$ be defined on the nonnegative integers by the recurrence

$$T(n) = aT(n/b) + f(n),$$

where we interpret n/b to mean either $\lfloor n/b \rfloor$ or $\lceil n/b \rceil$. Then $T(n)$ can be bounded asymptotically as follows.

1. If $f(n) = O(n^{\log_b a - \epsilon})$ for some constant $\epsilon > 0$, then $T(n) = \Theta(n^{\log_b a})$.

2. If $f(n) = \Theta(n^{\log_b a})$ then $T(n) = \Theta(n^{\log_b a} \lg n)$.

3. If $f(n) = \Omega(n^{\log_b a + \epsilon})$ for some constant $\epsilon > 0$, and if $af(n/b) \leq cf(n)$ for some constant $c < 1$ and all sufficiently large n , then $T(n) = \Theta(f(n))$.

3. Exponential Identities. For all real $a \neq 0$, m and n :

$$a^0 = 1. \quad a^1 = a. \quad a^{-1} = 1/a. \quad (a^m)^n = a^{mn}.$$

$$(a^m)^n = (a^n)^m. \quad a^m a^n = a^{m+n}. \quad a^m / a^n = a^{m-n}.$$

4. Logarithmic Identities. For all $a > 0$, $b > 0$, $c > 0$ and n :

$$a = b^{\log_b a}. \quad \log_c(ab) = (\log_c a) + (\log_c b). \quad \log_b a^n = n \log_b a. \quad \log_b a = \frac{\log_c a}{\log_c b}.$$

$$\log_b(1/a) = -\log_b a. \quad \log_b a = 1/\log_a b. \quad a^{\log_b n} = n^{\log_b a}. \quad n^{1/\log_2 n} = 2.$$

5. Some summations.

$$\sum_{k=1}^n k = \frac{1}{2}n(n+1). \quad \sum_{k=0}^n x^k = \frac{x^{n+1} - 1}{x - 1}. \text{ for real } x \neq 1$$

$$\sum_{k=0}^{\infty} x^k = \frac{1}{1-x}. \text{ when } |x| < 1. \quad \sum_{k=0}^{\infty} kx^k = \frac{x}{(1-x)^2}. \text{ when } |x| < 1.$$

$$\sum_{k=1}^n 1/k = \ln n + O(1). \quad \sum_{k=1}^{n-1} k \lg k \leq \frac{1}{2}n^2 \lg n - \frac{1}{8}n^2 \leq n^2 \lg n.$$